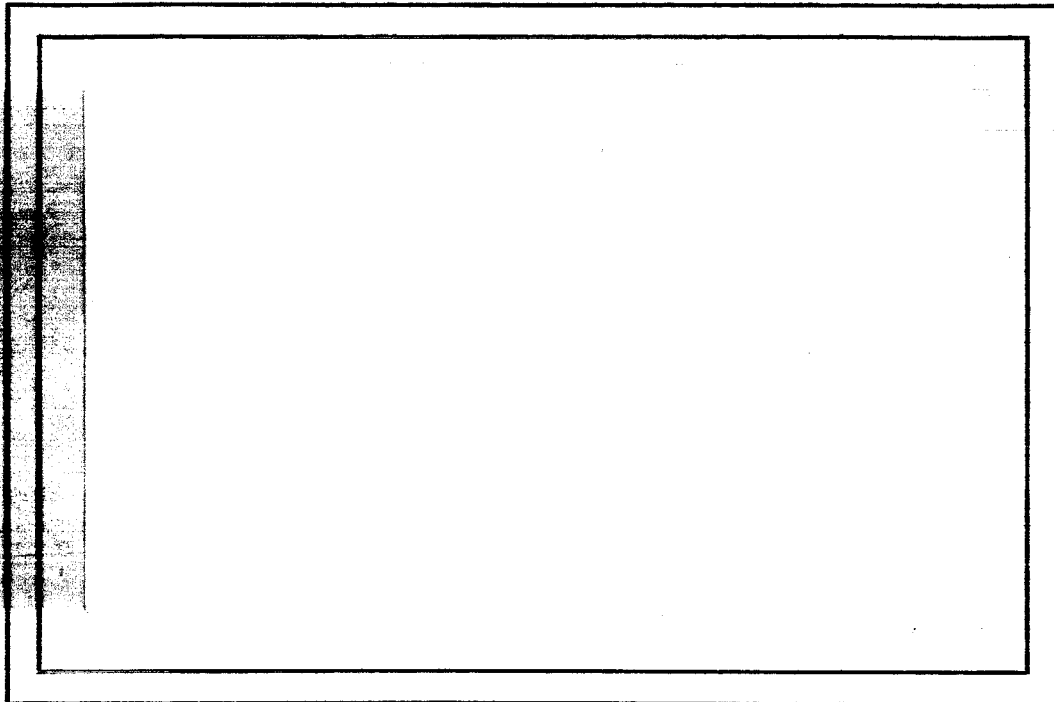


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*t*: The Use of Digital Computers to  
Determine Definitions for Abstract  
Groups

by

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N.A.S.A. - Trainee  
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ABSTRACT

This paper describes the work done by the author and others on digital computer programs for automatic enumeration of cosets. It is followed by a brief description of some of the work done on the finite "Burnside" groups especially by computer enumerations. A definition for the Burnside group  $B_{3,4}$  of exponent 3 with 4 generators involving only 35 words is given.

AUTHOR

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## CHAPTER I

### BASIC PROBLEM

If it is known that a group,  $G$ , is generated by a set of generators:  $S_1, S_2, \dots, S_k$  and that these generators satisfy a set of relations of the form:

$$f_i(S_1, S_2, \dots, S_k) = E, \quad (i = 1, 2, \dots, n) \quad (1)$$

then a natural question which arises is: What is the order of the group thus defined? An equivalent way of expressing this question is to ask how many of the combinations and permutations of the generators are distinct when subjected to the constraints implied by the relations,  $f_i$ , and the definition of an abstract group. For all but the most trivial of groups the direct approach clearly involves too much manipulation to be practical.

The most obvious technique to reduce the magnitude of the problem is to choose as large a subgroup,  $H$ , of  $G$  as possible -- the order of  $H$  being known from the set of defining relations -- and enumerate all of the left<sup>1</sup> cosets of that subgroup. For example, suppose one wished to determine the

---

<sup>1</sup>The choice of left or right cosets is immaterial to the process. The reason for choosing left cosets is explained in the paragraph of Chapter II.



order of the group,  $(8,7|2,3)^2$ , defined by the relations:

$$A^8 = B^7 = (AB)^2 = (A^{-1}B)^3 = E \quad (2)$$

There are several choices for a subgroup. One may use the group generated by B, denoted  $\{B\}$ , which is the cyclic group of order 7, or  $\{A\}$ , the cyclic group of order 8. However, a still better choice is  $K \equiv \{A^2, A^{-1}B\}$ .  $A^2$  is of period 4,  $A^{-1}B$  is of period 3, and their product,  $A^2 \cdot A^{-1}B = AB$ , is of period 2. These generators satisfy the relations (2,3,4):

$$R^2 = S^3 = (RS)^4 = E \quad (3)$$

which define the symmetric group on four elements,  $S_4$ , of order 24. This is easily seen by setting  $B^{-1}A = S$  and  $AB = R$ . Then  $RS = AB \cdot B^{-1}A = A^2$ . The order of  $(8,7|2,3)$  is known to be 10,752; therefore, using the subgroups specified only 1536, 1344, and 448 cosets would have to be enumerated respectively according as the basic subgroup is  $\{B\}$ ,  $\{A\}$  or K.

It is well known from elementary group theory that distinct cosets of the subgroup, H, have empty intersections; that any given coset, aH, is independent of the choice of the

---

<sup>2</sup>This group, which was first described by Dr. A. Sinkov: "On the group-Defining Relations (2,3,7;p)" Annals of Mathematics XXXVIII (1937), pp 580-582 and was later given an irreducible definition, has been widely used as a test problem for computer enumeration problems and thus for a comparison of the various techniques.

representative,  $a$ , that is  $b \in aH \Rightarrow aH = bH$ ; and that the union of the collection of all cosets of  $H$  is the entire group,  $G$ :

$$G = a_1 H \cup a_2 H \cup \dots \cup a_m H \quad (4)$$

where  $a_i \in H$  and  $a_i H \cap a_j H = \emptyset$ ,  $i \neq j$

Thus if in an enumeration of  $G$  by cosets of  $H$ ,  $m$  cosets were defined, the order of  $G$  would be  $m$  times the order of  $H$ .

## CHAPTER II

### SYSTEMATIC ENUMERATION

In 1936 Todd and Coxeter<sup>3</sup> presented a mechanical method for accomplishing this type of enumeration. Except for changes involving the choice of considering the elements of the group,  $G$ , as right multipliers, the process to be described below is that of Todd and Coxeter. The choice of right multipliers (left cosets) produces tables which are read in the conventional manner.

Consider a typical member of the set of relations (1).

It would have the form:

$$R_1 R_2 \dots R_l = E \quad (5)$$

where each  $R_i$  is one of the generators or the inverse of one of the generators and generators are repeated if necessary.

Such a relation is called a word of  $l$  letters. For example, if the commutator of generators  $A$  and  $B$  were to have its period specified as 2, (5) would have the form:

$$A^{-1} B^{-1} A B A^{-1} B^{-1} A B = E \quad (6)$$

In the process about to be described, there is a coset multiplication table for each word in the set of defining

---

<sup>3</sup>J. A. Todd and H. S. M. Coxeter. "A Practical Method for Enumerating Cosets of a Finite Abstract Group," Proceedings of the Edinburgh Mathematical Society, V, Series II (1936), pp. 26-34.

relations, (1). Each table has one column more than the number of letters in the corresponding word. The generators forming the word are placed at the head of the table between the various columns as in the following example:  $S_4$  defined by the relations (3).

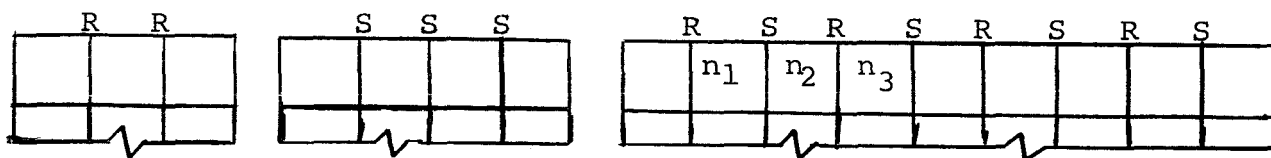


Fig. 1 -- Diagram of typical enumeration table

In Figure 1 the numbers  $n_1$ ,  $n_2$ , and  $n_3$  represent cosets. The entries are to be interpreted as  $n_1 \cdot S = n_2$  and  $n_2 \cdot R = n_3$  and, when reading backwards, as  $n_3 \cdot R^{-1} = n_2$  and  $n_2 \cdot S^{-1} = n_1$ . The constraints implied by the defining relations (1) are applied by requiring the coset numbers at both ends of a row to be identical.

The process is initiated by inserting in the tables the choice for coset 1, namely the subgroup  $H$ . In the example cited above, a reasonable choice for  $H$  is the subgroup generated by  $S$ , the cyclic group of order 3. Hence we denote  $H = S = 1$ . It follows that  $1 \cdot S = 1$  since any element  $H$  is of the form  $S$  and  $S \cdot S$  is also of the form  $S$ . Thus the above result is entered into the tables in every possible position as follows:



In Figure 3 it is deduced that  $2 \cdot R = 1$ . Algebraically this deduction is a consequence of the relation  $R^2 = E$ . Coset 2 is of the form  $S^\alpha R$ . Thus  $2 \cdot R$  is of the form  $S^\alpha \cdot R^2 = S^\alpha E = S^\alpha = 1$ . As the process continues, more instances of deduced information will occur. The advantage of this mechanical procedure is that the algebraic consequences of the defining relations are handled automatically, thus eliminating the necessity for abstract manipulation.

Also note that  $(RS)^4 = E$  is a relation of the form (7) as in fact are the other two. Therefore, a row beginning and ending with coset 2 is superfluous and hence was omitted in Figure 3.

The next pair of definitions might be  $2 \cdot S = 3$  followed by  $3 \cdot S = 4$  from which it is deduced that  $4 \cdot S = 2$ . At this point, the tables appear as follows:

R R			S S S				R S R S R S R S								
1	2	1	1	1	1	1	1	2	3			4	2	1	1
3		3	2	3	4	2	4							3	4
4		4													

Fig. 4.--Later stage of enumeration

From this point one might proceed as follows. Define  $3 \cdot R = 5$  thus deducing that  $5 \cdot R = 3$ . Then define  $4 \cdot R = 6$  deducing that  $6 \cdot R = 4$  and  $5 \cdot S = 6$ . At this point the tables appear as follows:

R R			S S S				R S R S R S R S								
1	2	1	1	1	1	1	1*	2	3*	5	6*	4	2*	1	1
3	5	3	2	3	4	2	4*	6					5*	3	4
4	6	4	5	6		5									

Fig. 5.--Still later stage of enumeration

Note that the starred coset numbers are in symmetrical positions and hence only two rows need to be carried in the chart corresponding to the relation  $(RS)^4 = E$ .

The next pair of definitions might then be  $6 \cdot S = 7$  which leads to  $7 \cdot S = 5$  and finally  $7 \cdot R = 8$  from which the results  $8 \cdot R = 7$  and  $8 \cdot S = 8$  are easily found. At this point the tables have "closed up" and the process is complete. The resulting tables appear below:

R R			S S S				R S R S R S R S								
1	2	1	1	1	1	1	1	2	3	5	6	4	2	1	1
3	5	3	2	3	4	2	4	6	7	8	8	7	5	3	4
4	6	4	5	6	7	5									
7	8	7	8	8	8	8									

Fig. 6.--Completed enumeration

There were 8 cosets defined and since  $H$  is of order 3, the order of the group  $(2,3,4)$  is 24. The coset multiplication table may be summarized in the following diagram:

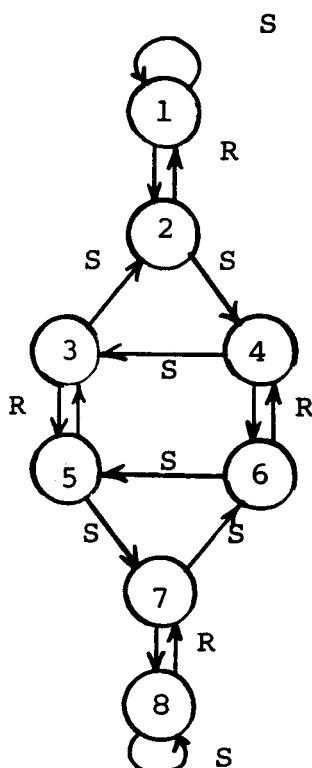


Fig. 7.--Coset diagram for  $(2,3,4)$

Any path from 1 to  $n$  yields an expression for an element of coset  $n$ . For example, the element  $SRSRS^2R$  in coset 8 corresponds to the path 1-1-2-4-6-5-7-8 on Figure 7.

In the process, it may occur that two different coset numbers are discovered to actually represent the same coset. When such a coincidence occurs, the larger number is replaced throughout the tables by the smaller and the process continues. When working by hand it may be convenient to use the larger number as that of the next coset defined.



As an example of a simple enumeration which involves a coincidence consider Klein's simple group of order 168 defined by the relations  $(2,3,7;4)$ :

$$P^2 = Q^3 = (QP)^7 = (Q^{-1}P^{-1}QP)^4 = E \quad (8)$$

Let  $PQ = R$  and  $Q^{-1} = S$ . Then we have  $P = PQ \cdot Q^{-1} = RS$  and  $Q^{-1}P^{-1}QP = (PQ)^{-1} Q \cdot PQ \cdot Q^{-1} = R^{-1}S^{-1}RS$ . Thus relations (8) imply

$$R^7 = S^3 = (RS)^2 = (R^{-1}S^{-1}RS)^4 = E \quad (9)$$

We now define the basic subgroup as  $\{R\}$  so that  $1 \cdot R = 1$  and proceed as indicated in the following table:

TABLE 1: DEFINITIONS FOR  $(2,3,7;4)$

Definition	Implied Consequences		
$1 \cdot S = 2$			
$2 \cdot S = 3$	$3 \cdot S = 1$	$2 \cdot R = 3$	
$3 \cdot R = 4$			
$4 \cdot S = 5$	$5 \cdot R = 2$		
$5 \cdot R^{-1} = 6$			
$6 \cdot S^{-1} = 7$			
$7 \cdot R = 8$			
$8 \cdot S = 9$			
$9 \cdot R^{-1} = 10$			
$10 \cdot S^{-1} = 11$			
$11 \cdot R = 12$	$12 \cdot S = 4$	$5 \cdot S = 12$	$12 \cdot R = 7$
$4 \cdot R = 13$	$13 \cdot S = 11$	$10 \cdot S = 13$	
$13 \cdot R = 14$	$14 \cdot R = 6$		
$11 \cdot R^{-1} = 15$	$9 \cdot S = 15$	$15 \cdot S = 8$	
$9 \cdot R = 16$	$16 \cdot S = 7$	$6 \cdot S = 16$	
$14 \cdot S^{-1} = 17$	$16 \cdot R = 17$		
$17 \cdot R = 18$			
$18 \cdot S = 19$			
$19 \cdot R^{-1} = 20$			
$20 \cdot S^{-1} = 21$			
$21 \cdot R = 22$			
$22 \cdot S = 23$			
$23 \cdot R^{-1} = 24$			

At this point the tables appear as in figure 8:

R R R R R R R							
1	1	1	1	1	1	1	1
2	3	4	13	14	6	5	2
7	8			15	11	12	7
9	16	17	18			10	9
19						20	19
21	22						21
23						24	23

S S S			
1	2	3	1
4	5	12	4
6	16	7	6
8	9	15	8
10	13	11	10
14		17	14
18	19		18
20		21	20
22	23		22
24			24

R S R S				
1	1	2	3	1
3	4	5	2	3
4	13	11	12	4
6	5	12	7	6
7	8	9	16	7
8			15	8
10	9	15	11	10
13	14		10	13
14	6	16	17	14
17	18	19		17
18				18
20	19		21	20
21	22	23		21
22				22
24	23			24

Fig. 8 -- (2,3,7;4) before the first coincidence

If the next definition made is:  $24 \cdot S^{-1} = 25$  we deduce from  $(R^{-1}S^{-1}RS)^4 = E$  that  $25 \cdot R = 10$  and then from  $(RS)^2 = E$  that  $14 \cdot S = 25$ . At this point we have coset 25 in two previously distinct rows in the table corresponding to  $S^3 = E$ . Comparing the rows we discover that  $17 \equiv 24$ , i.e. that the cosets represented by these two numbers are really the same. Next by comparing the rows containing 17 and 24 in the tables corresponding to  $R^7 = E$  we deduce that  $18 \equiv 23$ . A further search of the tables quickly indicates that no further consequential coincidences result. Thus we replace 24 by 17, 23 by 18, and 25 (the only coset defined after 23 and 24) by 23 (the lowest vacated coset number) and eliminate the rows which are duplicated. The resulting tables are then shown in figure 9:

1	1	1	1	1	1	1	1
2	3	4	13	14	6	5	2
7	8			15	11	12	7
9	16	17	18		23	10	9
19						20	19
21	22						21

1	2	3	1
4	5	12	4
6	16	7	6
8	9	15	8
10	13	11	10
14	23	17	14
18	19	22	18
20		21	20

1	1	2	3	1
3	4	5	2	3
4	13	11	12	4
6	5	12	7	6
7	8	9	16	7
8			15	14
10	9	15	11	10
13	14	23	10	13
14	6	16	17	14
17	18	19		17
20	19		21	20
21	22	18		21
22				22

1	1	3	4	5	6	7	8	9	10	11	12	4	3	2	3	1
2	5	4	13	11	15	9	16	7	12	5	2	3	2	1	1	2
6	14	17	18	19	20	21	22	18	17	23	10	13	4	12	7	6
8	7	16	17	14	13	10	9	15							15	8
10	9	8											8	15	11	10
12	11	13	14									17	16	6	5	12
16	9	8												14	6	16
20															21	20
21																21
22	21															22

Fig. 9--(2,3,7;4) after processing the first coincidence

We may then proceed as before to finish the enumeration. It turns out with one scheme of definitions (that used by the author's computer enumeration program) that 27 cosets will have to be defined and then a coincidence will reduce the number to 24 at which point the tables close-up. Thus the group  $(2,3,7;4)$  is of order  $7 \cdot 24 = 168$  as expected.

The order of defining new cosets is completely immaterial to the success of an enumeration. However, by a judicious sequence of definitions, the number of coincidences may be minimized. Experience seems to indicate that many groups cannot be enumerated without the occurrence of coincidences. As interesting examples of this I refer to John Leech's recent paper<sup>4</sup> where he cites private correspondence with Todd suggesting that two groups: Klein's simple group of order 168 defined by:

$$\begin{aligned} B^7 &= (AB)^2 = (A^{-1}B)^3 = (A^2B^2)^4 = E \\ \text{or} \\ B^7 &= (AB)^2 = (A^{-1}B)^3 = (A^3B^4)^3 = E \end{aligned} \tag{10}$$

and the previously cited group,  $(8,7|2,3)$ , defined by (2), when enumerated as cosets of  $\{B\}$  and  $\{A^2, A^{-1}B\}$  respectively, cannot be enumerated without the occurrence of coincidences.

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<sup>4</sup>John Leech. "Coset Enumeration on Digital Computers," Proceedings of the Cambridge Philosophical Society, LIX (1963), 285

# CHAPTER III

## ENUMERATION BY MACHINE

### Introduction

In 1957 Coxeter and Moser stated in the introduction to their book that the "method (for systematic enumeration) is sufficiently mechanical for the use of an electronic computer."<sup>5</sup> Since then several people have independently written programs for various machines to accomplish this task.

Leech gave a history of the work done on this problem that was known to him at the time of publication of his paper.<sup>6</sup> He gave an excellent description of the work of C. B. Haselgrove and his own work. He then cited the work of R. Maddison and A. Sinkov. All of the above mentioned programs used basically the same logic and of these, the work of Sinkov is best known to the author of this paper.

### Description of Logic Used by Sinkov

The first important way in which all of the computer programs differ from the hand method is in the elimination

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<sup>5</sup> H. M. S. Coxeter and W. O. T. Moser, Generators & Relations for Discrete Groups (Berlin: Springer-Verlag 1957), p. v.

<sup>6</sup> Leech, pp. 259-263

of the tables for each relation. Instead, tables are carried only for each generator and its inverse. This results in a considerable saving of storage space and therefore permits a larger group to be enumerated. The defining relations are stored in the machine and the generators comprising the relations are fetched as needed.

Sinkov's program essentially applies each coset in turn to each of the defining relations in turn. Assume that  $a_0$  is the current coset number, that  $m$  cosets have already been defined, and that  $f_i(s_1, s_2, \dots, s_n) = R_1 R_2 \dots R_l = E$  is the current defining relation. Then  $a_0 R_1$  is extracted from the table for the generator  $R_1$ . If  $a_0 R_1$  is defined e.g.  $a_0 R_1 = a_1$ , then  $a_1 R_2$  is extracted from the table and so on. If for some  $j$ ,  $a_j R_{j+1}$  is not defined, it is immediately defined as coset  $m+1$ , the appropriate entries are made in the tables, and the processing continues. When the end of the relation is reached, a test is made to determine if  $a_0 = a_l$ . If  $a_0 \neq a_l$ , a coincidence has been discovered. If  $a_0 = a_l$  or if a coincidence and all consequential coincidences have been processed, the current coset,  $a_0$ , is applied in like manner to the next relation. When the coset  $a_0$  has been applied to all of the relations, the next coset  $a_0+1$  is applied to the relations in turn. The process is complete when the last coset defined has been applied to all of the relations, without causing any new definitions to be introduced.



When a coincidence,  $a \equiv b$ , with  $a < b$ , is discovered, the row corresponding to coset  $b$  is examined for all generators and inverses. If a given entry,  $bR_i$ , is undefined, no action is necessary and the next entry is examined. Otherwise a test is made to see if  $bR_i = b$  and if that is the case, it is replaced by  $a$ . If  $bR_i \neq b$ , then the inverse entry  $(bR_i)R_i^{-1}$  is deleted from the table. Next the entry  $aR_i$  is examined. If  $aR_i$  is not defined, the entry  $bR_i$  is inserted. If  $aR_i$  is defined and  $aR_i = b$  it is replaced by  $a$ . Otherwise a new coincidence is set up between  $aR_i$  and  $bR_i$ . Then a check is made to determine if  $(aR_i)R_i^{-1}$  is defined and if not,  $a$  is inserted. Finally the entries in the row  $b$  are deleted (made zero).

The list of coincidences awaiting processing is sorted lexicographically so that redundant information need not be stored and also to assure that no coincidence is processed on a row already made zero.

After the entire list of coincidences has been processed, it is desirable for efficient use of memory space to eliminate the vacated rows from all the tables. This is easily done by using the coincidence routine to set up an artificial coincidence between the first empty row and the next non-empty one. This process is repeated until the tables are again without empty rows.

Description of Logic Used by Author

The program written by the author of this paper uses the other logic scheme presented in the literature. The method is essentially that of H. Felsch<sup>7</sup> although the program was written independently. This method was also used by Bandler (see Leech's paper<sup>8</sup>) although he did not program for automatic processing of coincidences.

The author's program was originally written for the IBM 1620, but in the spring of 1963 an IBM 7090 was delivered to the Computer Center of the University of Maryland, so the program was rewritten and modified using Fortran II for the 7090. Fortran II is a problem oriented programming language and hence a source listing of the program (see Appendix) may be of interest.

Basically the procedure used in this program is as follows: The cosets are applied sequentially to the defining relations in turn as before; however, when the forward working is halted by an undefined coset, the current coset is then applied to the inverse of the last generator in the relation. This backward working proceeds in a manner similar to the forward working and one of three things may happen. First,

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<sup>7</sup>H. Felsch, "Programmierung der Restklassenabzaehlung einer Gruppe nach Untergruppen," Numerische Mathematik, III (1961), pp. 250-256

<sup>8</sup>Leech, p. 262.

the backward working may encounter an undefined coset in which case a new coset is defined; second, the backward working may just meet the forward working, in which case new information has been deduced; or third, the forward working and the backward working may overlap, in which case a coincidence is deduced. When a new coset is defined or when new information is deduced, e.g.,  $aR_i = b$ , every occurrence of the generator  $R_i$  or its inverse  $R_i^{-1}$  is examined. The coset  $a$  is applied to the word in which the generator  $R_i$  appeared shifted cyclically to begin with  $R_i$ . The same procedure as in the general working is used, namely, upon reaching a gap in the forward working, backward working is begun. However, if a void is discovered a new definition is not made. New information discovered in this manner is entered into the multiplication tables and stored away for future processing. Coincidences are handled in the same manner as in Sinkov's program except that provision is made to make any necessary changes to the table of information awaiting processing that might have occurred due to the processing of subsequent coincidences.

In Leech's paper he stated, "No direct comparison of running times with the two methods is available at present as the machine speeds are widely different; this must wait

until both methods have been programmed for the same machine."<sup>9</sup>  
 A partial answer to this question is now available since Sinkov's problem was written for the IBM 704 and is capable of being run on the IBM 7090. The following chart from a report written by Sinkov for the Computer Science Center at the University of Maryland in June 1963 shows a comparison of running times on the classical problem  $(8,7|2,3)$ . The author's running time has been added.

TABLE I  
 COMPARISON OF COMPUTER RUNNING TIMES  
 ON THE GROUP  $(8,7|2,3)$

Person	Machine	Cosets Required	Time
Todd	By hand	945	>30 hours
Felsch	Zeus 22	1300	~2 hours
Leech	EDSAC 2	2000	42 minutes
Sinkov	IBM 7090	2176	5 minutes
Leech	KDF 9		2 min. 30 sec.
Snively	IBM 7090	1747	1 min. 36 sec.

The faster time obtained by the author's program is not necessarily indicative of more efficient logic. The author's logic is considerably more complicated and therefore takes up more storage space thus limiting the size of the problem that may be handled within the memory of the machine. In

<sup>9</sup>Leech, p. 263

fact the author's program can enumerate  $10000/n$  cosets where  $n$  is the number of generators of the group being enumerated. Running times are expected to vary widely from problem to problem. The Felsch logic would be better in a problem that involves a large number of excess cosets to be defined since relatively few excess cosets are defined by the author's scheme.

## CHAPTER IV

### THE BURNSIDE PROBLEM

In 1902 Burnside stated what is commonly called the Burnside problem: can the order of a group with a finite number of generators be infinite while the period of each element in the group is finite?<sup>10</sup> To the author's knowledge, the problem remains unsolved, for although a Russian mathematician, Novikov, claimed to have answered it affirmatively, his proof has not yet been published.

A more specialized problem may be stated simply: assume the groups under consideration are finitely generated and that the orders of every element in the group, are bounded. Suppose, for example,  $S_1, S_2, \dots, S_r$  generate a group,  $B_{n,r}$ , and every element  $R \in B_{n,r}$  satisfies the relation  $R^n = E$ . Then  $B_{n,r}$  is called the Burnside group of exponent  $n$  with  $r$  generators. This Burnside problem now reduces to the question: which of the groups  $B_{n,r}$  are finite?

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<sup>10</sup>W. Burnside, "On an Unsettled Question in the Theory of Discontinuous Groups," Quarterly Journal of Pure and Applied Mathematics, XXXIII (1902), p. 230-238

In his book Marshall Hall summarizes the work done on this problem up to 1959.<sup>11</sup> Some interesting results are:

1.  $B_{2,r}$  is finite (for finite  $r$ ), namely the abelian group of the form:

$$\begin{array}{c} C_2 \times C_2 \times \dots \times C_2 \\ r \text{ factors} \end{array} \quad (11)$$

2.  $B_{3,r}$  is finite and of order:

$$\frac{r(r^2+5)}{6} \quad (12)$$

This result was obtained by Levi and van der Waerden.<sup>12</sup>

3.  $B_{4,r}$  is finite. This result was obtained by Sanov.<sup>13</sup>
4.  $B_{5,2}$  if it is finite, has order at most  $5^3$  (see Kostrikin)<sup>14</sup>

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<sup>11</sup>M. Hall, The Theory of Groups. (New York: The MacMillan Company, 1959), pp. 320-338

<sup>12</sup> Ibid., p. 321

<sup>13</sup> Ibid., p. 324

<sup>14</sup> Ibid., p. 327

5.  $B_{6,r}$  is finite and of order:

$$\frac{b(b^2+5)}{6} \\ 2^a 3$$

where:

$$a = 1 + (r-1)3 \frac{r(r^2+5)}{6} \quad (13)$$

and:

$$b = 1 + (r-1)2^r$$

This result was obtained by M. Hall.<sup>15</sup>

Although the orders of the groups  $B_{3,r}$  are known, sets of irreducible or nearly irreducible defining relations are not known for all of these groups. In attempting to find such sets of defining relations, one finds a good application for computer enumeration. The technique used is to overdefine the group, that is, to fix the periods of more elements of the group than is necessary to define the group. Then when the order of the group is thus determined (if not already known) the enumeration is rerun with some of the relations removed. If a group of the same order results, the relations removed were redundant, i.e., an algebraic consequence of the remaining relations. By proceeding in this manner, a non-redundant

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<sup>15</sup> Ibid., p. 336-337



set of relations may be obtained.

$B_{3,1}$  is, of course, the cyclic group of order 3 satisfying the relation:

$$A^3 = E \quad (14)$$

$B_{3,2}$ , of order 27, satisfies the following set of relations:

$$A^3 = B^3 = (AB)^3 = (A^{-1}B)^3 = E \quad (15)$$

In his 1963 paper, Leech gave his results for the group  $B_{3,3}$ .<sup>16</sup> He obtained his results by enumerating the 81 cosets of  $A, B$  which is  $B_{3,2}$  of order 27. One of the resulting definitions is:

$$\begin{aligned} A^3 = B^3 = C^3 = (AB)^3 = (AC)^3 = (BC)^3 = (A^{-1}B)^3 = \\ (A^{-1}C)^3 = (B^{-1}C)^3 = (ABC)^3 = (A^{-1}BC)^3 = (AB^{-1}C)^3 = \\ (ABC^{-1})^3 = E \end{aligned}$$

His enumeration was performed on EDSAC 2. The memory of EDSAC 2 was not large enough to permit the enumeration of  $B_{3,4}$ . This is the problem solved by the author on the IBM 7090.

In approaching the problem of  $B_{3,4}$  with generators  $A, B, C, D$  the first step was to consider the generators in all combinations of three and assure that they satisfy the relations (14). When this is done, we are assured that all words containing only three of the four generators are of exponent 3.

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<sup>16</sup>Leech, p. 264

To accomplish this, it was sufficient to specify the periods of 32 elements. These are:

$$\begin{aligned}
 A^3 &= B^3 = C^3 = D^3 = E \\
 (AB)^3 &= (AC)^3 = (AD)^3 = (BC)^3 = (BD)^3 = (CD)^3 = E \\
 (A^{-1}B)^3 &= (A^{-1}C)^3 = (A^{-1}D)^3 = (B^{-1}C)^3 = (B^{-1}D)^3 = (C^{-1}D)^3 = E \\
 (ABC)^3 &= (A^{-1}BC)^3 = (AB^{-1}C)^3 = (ABC^{-1})^3 = E \\
 (ACD)^3 &= (A^{-1}CD)^3 = (AC^{-1}D)^3 = (ACD^{-1})^3 = E \\
 (BCD)^3 &= (B^{-1}CD)^3 = (BC^{-1}D)^3 = (BCD^{-1})^3 = E
 \end{aligned} \tag{17}$$

Let  $W_{ij} = R_1^{\xi_1} R_2^{\xi_2} R_3^{\xi_3} R_4^{\xi_4}$  where  $(i = 1, 2, \dots, 6)$  and  $(j = 1, 2, \dots, 8)$  such that if  $Y_j = (\xi_1, \xi_2, \xi_3, \xi_4)$  we have

$$\begin{aligned}
 Y_1 &= (1, 1, 1, 1) & Y_5 &= (1, 1, 1, -1) \\
 Y_2 &= (-1, 1, 1, 1) & Y_6 &= (-1, -1, 1, 1) \\
 Y_3 &= (1, -1, 1, 1) & Y_7 &= (-1, 1, -1, 1) \\
 Y_4 &= (1, 1, -1, 1) & Y_8 &= (-1, 1, 1, -1)
 \end{aligned} \tag{18}$$

and

$$\begin{aligned}
 W_{11} &= ABCD & W_{41} &= ACDB \\
 W_{21} &= ABDC & W_{51} &= ACBC \\
 W_{31} &= ACBD & W_{61} &= ADCB
 \end{aligned} \tag{19}$$

It is readily verified that if A, B, C, and D satisfy relations (17) and

$$W_{ij}^3 = E \quad (i = 1, 2, \dots, 6; j = 1, 2, \dots, 8) \tag{20}$$



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Since an enumeration of relations (17) alone failed to give closure before memory capacity was exceeded it is quite likely that the last eight successful sets of defining relations in (22) are irreducible. A complete proof of the last statement by enumerations would require 35 enumerations for each definition or about 72 hours of 7090 computer time.

## CHAPTER V

### OUTLOOK FOR THE FUTURE

The obvious next step would be to attempt to determine a set of defining relations for  $B_{3,5}$ . This, however, is a problem which far exceeds the capability of the 7090 since  $B_{3,5}$  is of order  $3^{25}$  and the largest subgroup available for an enumeration is  $B_{3,4}$  of order  $3^{14}$  which would require a total of  $3^{11}$  cosets to be defined. Since there are tables for each generator and its inverse, a total of 10 tables, this means that a total of 1,771,470 table entries must be provided. Even packing two entries to a word ( $3^{11} < 2^{18}$ ) only 65,536 entries can be provided, not allowing room for the program and other tables.

A means of extending the program's capabilities would be to store the tables on magnetic tape and call them into memory as needed; however, this is very impractical because such operations are quite time consuming and large amounts of computer time are not readily available.

Another possibility for solving this problem is a disk storage similar to the IBM 1301, but unfortunately this was not available to the author.

However, the work on  $B_{3,4}$  did permit a conjecture. Given the definition for group  $B_{3,n}$ , it seems likely that in addition to the combinations of relations needed to define

$B_{3,n}$  as a subgroup of  $B_{3,n+1}$  in  $\frac{(n+1)n}{2}$  ways it is only necessary to add a portion of the words of  $n+1$  letters to completely define  $B_{3,n+1}$ .

Another interesting question is the study of  $B_{5,2}$ ; however, the largest readily known subgroup available is  $B_{5,1} = C_5$ , the cyclic group of order 5. Even if the order of  $B_{5,2}$  were as low as  $5^{10}$  one would have to enumerate  $5^9 = 1,953,125$  cosets, a task which is well beyond the capability of the 7090 without an extremely large random access storage.

In the future one may expect computers to become faster and to have larger memories. At the current machine speed a memory of 5,000,000 IBM 7090 words would enable the author's program to undertake the problem of  $B_{3,5}$  and possibly  $B_{5,2}$ , however, the time required to run these problems would be prohibitive. With a hundred fold increase in processing speed these two problems would be well within the range of machine enumeration.

APPENDIX

SOURCE LISTING OF  
COSET ENUMERATION PROGRAM

		MAIN	
*	LIST8	MAIN	01
*	LABEL	MAIN	02
C	7090 COSET EVALUATION - FELSCH LOGIC - JAMES W. SNIVELY, JR.	MAIN	03
	DIMENSION ICHRT(20000),IC0IN1(1000),IC0IN2(1000),IGEN(100,13)	MAIN	04
	DIMENSION IPOWR(100),NLETR(100),ITBL(10,200),IWPG(10),ITALLY(100)	MAIN	05
	DIMENSION INFO1(50),INFO2(50),INFO3(50)	MAIN	06
	COMMON NWORD,NORD,NUM,NGENS,MAX,MAXROW,NDEF,IWORK,IWAIT,IR	MAIN	07
	COMMON LPWMAX,LWPGMX,ICTMAX,IWMAX,ICMAX,IN,NOUT,ITR,NJOB,NJOBS	MAIN	08
	COMMON ISCAN,ILOOK,ITALLY,INFO1,INFO2,INFO3,IWPG,ITBL	MAIN	09
	COMMON NLETR,IPOWR,IGEN,IC0IN1,IC0IN2,ICHRT	MAIN	10
	NJOB=1	MAIN	11
	IN=5	MAIN	12
	NOUT=6	MAIN	13
	LPWMAX=14	MAIN	14
	LWPGMX=200	MAIN	15
	ICTMAX=20000	MAIN	16
	IWMAX=50	MAIN	17
	ICMAX=1000	MAIN	18
	IWAIT=0	MAIN	19
	READ INPUT TAPE IN,60,NJOBS	MAIN	20
10	CALL CLOCK(6H= TIME,-1)	MAIN	21
	CALL INPUT	MAIN	22
	CALL ARRAY	MAIN	23
	CALL RESET	MAIN	24
	CALL SCAN	MAIN	25
	IF(NJOB-NJOBS) 40,50,50	MAIN	26
40	NJOB=NJOB+1	MAIN	27
	GO TO 10	MAIN	28
50	CALL CLOCK(6H= TIME,-1)	MAIN	29
	CALL EXIT	MAIN	30
60	FORMAT(I5)	MAIN	31
	END	MAIN	32
		INPUT	
*	LIST8	INPUT	01
*	LABEL	INPUT	02
	SUBROUTINE INPUT	INPUT	03
	DIMENSION ICHRT(20000),IC0IN1(1000),IC0IN2(1000),IGEN(100,13)	INPUT	04
	DIMENSION IPOWR(100),NLETR(100),ITBL(10,200),IWPG(10),ITALLY(100)	INPUT	05
	DIMENSION INFO1(50),INFO2(50),INFO3(50)	INPUT	06
	COMMON NWORD,NORD,NUM,NGENS,MAX,MAXROW,NDEF,IWORK,IWAIT,IR	INPUT	07
	COMMON LPWMAX,LWPGMX,ICTMAX,IWMAX,ICMAX,IN,NOUT,ITR,NJOB,NJOBS	INPUT	08
	COMMON ISCAN,ILOOK,ITALLY,INFO1,INFO2,INFO3,IWPG,ITBL	INPUT	09
	COMMON NLETR,IPOWR,IGEN,IC0IN1,IC0IN2,ICHRT	INPUT	10
	READ INPUT TAPE IN,20,NWORD	INPUT	11
	DO 10 I=1,NWORD	INPUT	12
	READ INPUT TAPE IN,30,NLETR(I),IPOWR(I)	INPUT	13
	K=NLETR(I)	INPUT	14
10	READ INPUT TAPE IN,40,(IGEN(I,J),J=1,K)	INPUT	15
	READ INPUT TAPE IN,50,NORD,ITR,NGENS	INPUT	16
	RETURN	INPUT	17
20	FORMAT(I5)	INPUT	18
30	FORMAT(2I5)	INPUT	19
40	FORMAT(13I5)	INPUT	20
50	FORMAT(3I5)	INPUT	21
	END	INPUT	22



* LIST8	ARRAY 01
* LABEL	ARRAY 02
SUBROUTINE ARRAY	ARRAY 03
DIMENSION ICHRT(20000),IC0IN1(1000),IC0IN2(1000),IGEN(100,13)	ARRAY 04
DIMENSION IPOWR(100),NLETR(100),ITBL(10,200),IWPG(10),ITALLY(100)	ARRAY 05
DIMENSION INFO1(50),INFO2(50),INFO3(50)	ARRAY 06
COMMON NWORD,NORD,NUM,NGENS,MAX,MAXROW,NDEF,IWORK,IWAIT,IR	ARRAY 07
COMMON LPWMAX,LWPGMX,ICTMAX,IWMAX,ICMAX,IN,NOUT,ITR,NJOB,NJOBS	ARRAY 08
COMMON ISCAN,ILOOK,ITALLY,INFO1,INFO2,INFO3,IWPG,ITBL	ARRAY 09
COMMON NLETR,IPOWR,IGEN,IC0IN1,IC0IN2,ICHRT	ARRAY 10
DO 40 I=1,NGENS	ARRAY 11
II=1	ARRAY 12
DO 30 J=1,NWORD	ARRAY 13
JJ=NLETR(J)	ARRAY 14
DO 30 K=1,JJ	ARRAY 15
IF(XABSF(IGEN(J,K))-1) 30,20,30	ARRAY 16
20 ITBL(I,II)=LPWMAX*J+K	ARRAY 17
II=II+1	ARRAY 18
IF(II-LWPGMX) 30,30,50	ARRAY 19
30 CONTINUE	ARRAY 20
40 IWPG(I)=II-1	ARRAY 21
RETURN	ARRAY 22
50 CALL ERROR(1)	ARRAY 23
RETURN	ARRAY 24
END	ARRAY 25
	RESET
* LIST8	RESET 01
* LABEL	RESET 02
SUBROUTINE RESET	RESET 03
DIMENSION ICHRT(20000),IC0IN1(1000),IC0IN2(1000),IGEN(100,13)	RESET 04
DIMENSION IPOWR(100),NLETR(100),ITBL(10,200),IWPG(10),ITALLY(100)	RESET 05
DIMENSION INFO1(50),INFO2(50),INFO3(50)	RESET 06
COMMON NWORD,NORD,NUM,NGENS,MAX,MAXROW,NDEF,IWORK,IWAIT,IR	RESET 07
COMMON LPWMAX,LWPGMX,ICTMAX,IWMAX,ICMAX,IN,NOUT,ITR,NJOB,NJOBS	RESET 08
COMMON ISCAN,ILOOK,ITALLY,INFO1,INFO2,INFO3,IWPG,ITBL	RESET 09
COMMON NLETR,IPOWR,IGEN,IC0IN1,IC0IN2,ICHRT	RESET 10
MAXROW=ICTMAX/(2*NGENS)	RESET 11
N=2*NGENS*MAXROW	RESET 12
DO 10 I=1,N	RESET 13
10 ICHRT(I)=0	RESET 14
READ INPUT TAPE IN,60,N,NDEF	RESET 15
IF(N) 40,40,20	RESET 16
20 DO 30 I=1,N	RESET 17
READ INPUT TAPE IN,70,J,K,L	RESET 18
CALL SEEK(-J,K,M)	RESET 19
ICHRT(M)=L	RESET 20
CALL SEEK(-L,-K,M)	RESET 21
30 ICHRT(M)=J	RESET 22
40 MAX=NDEF	RESET 23
IR=0	RESET 24
DO 50 I=1,NWORD	RESET 25
50 ITALLY(I)=0	RESET 26
RETURN	RESET 27

60	FORMAT(2I5)	RESET	28
70	FORMAT(3I5)	RESET	29
	END	RESET	30
		SCAN	
*	LIST8	SCAN	01
*	LABEL	SCAN	02
	SUBROUTINE SCAN	SCAN	03
	DIMENSION ICHRT(20000),IC0IN1(1000),IC0IN2(1000),IGEN(100,13)	SCAN	04
	DIMENSION IPOWR(100),NLETR(100),ITBL(10,200),IWPG(10),ITALLY(100)	SCAN	05
	DIMENSION INFO1(50),INFO2(50),INFO3(50)	SCAN	06
	COMMON NWORD,NORD,NUM,NGENS,MAX,MAXROW,NDEF,IWORK,IWAIT,IR	SCAN	07
	COMMON LPWMAX,LWPGMX,ICTMAX,IWMAX,ICMAX,IN,NOUT,ITR,NJOB,NJOBS	SCAN	08
	COMMON ISCAN,ILOOK,ITALLY,INFO1,INFO2,INFO3,IWPG,ITBL	SCAN	09
	COMMON NLETR,IPOWR,IGEN,IC0IN1,IC0IN2,ICHRT	SCAN	10
	GO TO (3000,1000,1000,1000),ITR	SCAN	11
1000	WRITE OUTPUT TAPE NOUT,2000	SCAN	12
2000	FORMAT(5HISCAN)	SCAN	13
3000	ISCAN=0	SCAN	14
10	ISCAN=ISCAN+1	SCAN	15
	DO 150 J=1,NWORD	SCAN	16
15	IWORK=ISCAN	SCAN	17
	KK=IPOWR(J)	SCAN	18
	LL=NLETR(J)	SCAN	19
	DO 20 K=1,KK	SCAN	20
	DO 20 L=1,LL	SCAN	21
	CALL SEEK(IWORK,IGEN(J,L),M)	SCAN	22
	GO TO (20,40),M	SCAN	23
20	CONTINUE	SCAN	24
	IF(IWORK-ISCAN) 30,150,30	SCAN	25
30	CALL COINC(IWORK,ISCAN,J)	SCAN	26
	GO TO 145	SCAN	27
40	IST=IWORK	SCAN	28
	IWORK=ISCAN	SCAN	29
	KT=K+1	SCAN	30
	LT=L+1	SCAN	31
	IF(KT-KK) 50,50,70	SCAN	32
50	DO 60 K=KT,KK	SCAN	33
	DO 60 N=1,LL	SCAN	34
	L=(LL+1)-N	SCAN	35
	CALL SEEK(IWORK,-IGEN(J,L),M)	SCAN	36
	GO TO (60,140),M	SCAN	37
60	CONTINUE	SCAN	38
70	IF(LT-LL) 80,80,100	SCAN	39
80	DO 90 N=LT,LL	SCAN	40
	L=(LL+LT)-N	SCAN	41
	CALL SEEK(IWORK,-IGEN(J,L),M)	SCAN	42
	GO TO (90,140),M	SCAN	43
90	CONTINUE	SCAN	44
100	L=LT-1	SCAN	45
	CALL SEEK(IWORK,-IGEN(J,L),M)	SCAN	46
	GO TO (110,120),M	SCAN	47
110	CALL COINC(IWORK,IST,J)	SCAN	48
	GO TO 145	SCAN	49
-120	CALL INFO(IST,IGEN(J,L),IWORK)	SCAN	50

GO TO 130	SCAN 51
130 CALL LOOKI(IST,IGEN(J,L),IWORK)	SCAN 52
GO TO 150	SCAN 53
140 L=LT-1	SCAN 54
CALL DEFINE(IST,IGEN(J,L))	SCAN 55
GO TO 15	SCAN 56
145 CALL CLOSE	SCAN 57
150 CONTINUE	SCAN 58
IF(ISCAN-NDEF) 10,160,160	SCAN 59
160 CALL FINISH	SCAN 60
RETURN	SCAN 61
END	SCAN 62
DEFINE	
* LIST8	DEFINE01
* LABEL	DEFINE02
SUBROUTINE DEFINE(INX,INY)	DEFINE03
DIMENSION ICHRT(20000),IC0IN1(1000),IC0IN2(1000),IGEN(100,13)	DEFINE04
DIMENSION IPOWR(100),NLETR(100),ITBL(10,200),IWPG(10),ITALLY(100)	DEFINE05
DIMENSION INFO1(50),INFO2(50),INFO3(50)	DEFINE06
COMMON NWORD,NORD,NUM,NGENS,MAX,MAXROW,NDEF,IWORK,IWAIT,IR	DEFINE07
COMMON LPWMAX,LWPGMX,ICTMAX,IWMAX,ICMAX,IN,NOUT,ITR,NJOB,NJOBS	DEFINE08
COMMON ISCAN,ILOOK,ITALLY,INFO1,INFO2,INFO3,IWPG,ITBL	DEFINE09
COMMON NLETR,IPOWR,IGEN,IC0IN1,IC0IN2,ICHRT	DEFINE10
GO TO (3000,1000,1000,1000),ITR	DEFINE11
1000 WRITE OUTPUT TAPE NOUT,2000,INX,INY	DEFINE12
2000 FORMAT(8H DEFINE(,I5,1H,,I5,1H))	DEFINE13
3000 IX=INX	DEFINE14
IY=INY	DEFINE15
NDEF=NDEF+1	DEFINE16
IF(NDEF-MAX) 20,20,10	DEFINE17
10 MAX=NDEF	DEFINE18
20 IF(NDEF-MAXROW) 40,40,30	DEFINE19
30 CALL ERROR(2)	DEFINE20
40 CALL SEEK(-IX,IY,1)	DEFINE21
ICHRT(1)=NDEF	DEFINE22
CALL SEEK(-NDEF,-IY,1)	DEFINE23
ICHRT(1)=IX	DEFINE24
CALL LOOKI(IX,IY,NDEF)	DEFINE25
RETURN	DEFINE26
END	DEFINE27
LOOKI	
* LIST8	LOOKI 01
* LABEL	LOOKI 02
SUBROUTINE LOOKI(INX,INY,INZ)	LOOKI 03
DIMENSION ICHRT(20000),IC0IN1(1000),IC0IN2(1000),IGEN(100,13)	LOOKI 04
DIMENSION IPOWR(100),NLETR(100),ITBL(10,200),IWPG(10),ITALLY(100)	LOOKI 05
DIMENSION INFO1(50),INFO2(50),INFO3(50)	LOOKI 06
COMMON NWORD,NORD,NUM,NGENS,MAX,MAXROW,NDEF,IWORK,IWAIT,IR	LOOKI 07
COMMON LPWMAX,LWPGMX,ICTMAX,IWMAX,ICMAX,IN,NOUT,ITR,NJOB,NJOBS	LOOKI 08
COMMON ISCAN,ILOOK,ITALLY,INFO1,INFO2,INFO3,IWPG,ITBL	LOOKI 09
COMMON NLETR,IPOWR,IGEN,IC0IN1,IC0IN2,ICHRT	LOOKI 10
GO TO (3000,3000,3000,1000),ITR	LOOKI 11
1000 WRITE OUTPUT TAPE NOUT,2000,INX,INY,INZ	LOOKI 12
2000 FORMAT(7H LOOKI(,I5,1H,,I5,1H,,I5,1H))	LOOKI 13

3000	IX=INX	LOOKI	14
	IY=INY	LOOKI	15
	IZ=INZ	LOOKI	16
	IWAIT=0	LOOKI	17
	CALL LOOK(IX,IY,IZ)	LOOKI	18
	IF(IWAIT) 40,40,20	LOOKI	19
20	I=IWAIT	LOOKI	20
	IWAIT=I-1	LOOKI	21
	CALL LOOK(INFO1(I),INFO2(I),INFO3(I))	LOOKI	22
	IF(IWAIT) 40,40,20	LOOKI	23
40	RETURN	LOOKI	24
	END	LOOKI	25
		LOOK	
*	LIST8	LOOK	01
*	LABEL	LOOK	02
	SUBROUTINE LOOK(INX,INY,INZ)	LOOK	03
	DIMENSION ICHRT(20000),IC0IN1(1000),IC0IN2(1000),IGEN(100,13)	LOOK	04
	DIMENSION IPOWR(100),NLETR(100),ITBL(10,200),IWPG(10),ITALLY(100)	LOOK	05
	DIMENSION INFO1(50),INFO2(50),INFO3(50)	LOOK	06
	COMMON NWORD,NORD,NUM,NGENS,MAX,MAXROW,NDEF,IWORK,IWAIT,IR	LOOK	07
	COMMON LPWMAX,LWPGMX,ICTMAX,IWMAX,ICMAX,IN,NOUT,ITR,NJOB,NJOBS	LOOK	08
	COMMON ISCAN,ILOOK,ITALLY,INFO1,INFO2,INFO3,IWPG,ITBL	LOOK	09
	COMMON NLETR,IPOWR,IGEN,IC0IN1,IC0IN2,ICHRT	LOOK	10
	GO TO (3000,3000,3000,1000),ITR	LOOK	11
1000	WRITE OUTPUT TAPE NOUT,2000,INX,INY,INZ	LOOK	12
2000	FORMAT(6H LOOK(,I5,1H,,I5,1H,,I5,1H))	LOOK	13
3000	IF(INY) 10,20,20	LOOK	14
10	ILOOK=INZ	LOOK	15
	IZ=INX	LOOK	16
	IY=-INY	LOOK	17
	GO TO 30	LOOK	18
20	ILOOK=INX	LOOK	19
	IZ=INZ	LOOK	20
	IY=INY	LOOK	21
30	II=IWPG(IY)	LOOK	22
	DO 230 I=1,II	LOOK	23
	IWORK=ILOOK	LOOK	24
	LL=ITBL(IY,I)/LPWMAX	LOOK	25
	KL=ITBL(IY,I)-LPWMAX*LL	LOOK	26
	JJ=IPOWR(LL)	LOOK	27
	KK=NLETR(LL)	LOOK	28
	DO 60 J=1,JJ	LOOK	29
	DO 60 M=1,KK	LOOK	30
	K=M+KL-1	LOOK	31
	IF(K-KK) 50,50,40	LOOK	32
40	K=K-KK	LOOK	33
50	CALL SEEK(IWORK,IGEN(LL,K),L)	LOOK	34
	GO TO (60,80),L	LOOK	35
60	CONTINUE	LOOK	36
	IF(IWORK-ILOOK) 70,230,70	LOOK	37
70	CALL COINC(IWORK,ILOOK,LL)	LOOK	38
	GO TO 225	LOOK	39
80	IST=IWORK	LOOK	40
	IWORK=ILOOK	LOOK	41

JT=J+1	LOOK	42
KT=M+1	LOOK	43
IF(JT-JJ) 90,90,130	LOOK	44
90 DO 120 J=JT,JJ	LOOK	45
DO 120 M=1,KK	LOOK	46
K=KK+KL-M	LOOK	47
IF(K-KK) 110,110,100	LOOK	48
100 K=K-KK	LOOK	49
110 CALL SEEK(IWORK,-IGEN(LL,K),L)	LOOK	50
GO TO (120,230),L	LOOK	51
120 CONTINUE	LOOK	52
130 IF(KT-KK) 140,140,180	LOOK	53
140 DO 170 M=KT,KK	LOOK	54
K=(KK+KL+KT)-(M+1)	LOOK	55
IF(K-KK) 160,160,150	LOOK	56
150 K=K-KK	LOOK	57
160 CALL SEEK(IWORK,-IGEN(LL,K),L)	LOOK	58
GO TO (170,230),L	LOOK	59
170 CONTINUE	LOOK	60
180 K=KL+KT-2	LOOK	61
IF(K-KK) 200,200,190	LOOK	62
190 K=K-KK	LOOK	63
200 CALL SEEK(IWORK,-IGEN(LL,K),L)	LOOK	64
GO TO (210,220),L	LOOK	65
210 CALL COINC(IWORK,IST,LL)	LOOK	66
GO TO 225	LOOK	67
220 CALL INFO(IST,IGEN(LL,K),IWORK)	LOOK	68
GO TO 230	LOOK	69
225 CALL CLOSE	LOOK	70
230 CONTINUE	LOOK	71
RETURN	LOOK	72
END	LOOK	73
	INFO	
* LIST8	INFO	01
* LABEL	INFO	02
SUBROUTINE INFO(IX,IY,IZ)	INFO	03
DIMENSION ICHRT(20000),IC0IN1(1000),IC0IN2(1000),IGEN(100,13)	INFO	04
DIMENSION IPOWR(100),NLETR(100),ITBL(10,200),IWPG(10),ITALLY(100)	INFO	05
DIMENSION INFO1(50),INFO2(50),INFO3(50)	INFO	06
COMMON NWORD,NORD,NUM,NGENS,MAX,MAXROW,NDEF,IWORK,IWAIT,IR	INFO	07
COMMON LPWMAX,LWPGMX,ICTMAX,IWMAX,ICMAX,IN,NOUT,ITR,NJOB,NJOBS	INFO	08
COMMON ISCAN,ILOOK,ITALLY,INFO1,INFO2,INFO3,IWPG,ITBL	INFO	09
COMMON NLETR,IPOWR,IGEN,IC0IN1,IC0IN2,ICHRT	INFO	10
GO TO (3000,1000,1000,1000),ITR	INFO	11
1000 WRITE OUTPUT TAPE NOUT,2000,IX,IY,IZ	INFO	12
2000 FORMAT(6H INFO(,I5,1H,,I5,1H,,I5,1H))	INFO	13
3000 CALL SEEK(-IX,IY,L)	INFO	14
ICHRT(L)=IZ	INFO	15
CALL SEEK(-IZ,-IY,L)	INFO	16
ICHRT(L)=IX	INFO	17
IWAIT=IWAIT+1	INFO	18
IF(IWAIT-IWMAX) 10,10,20	INFO	19
10 INFO1(IWAIT)=IX	INFO	20
INFO2(IWAIT)=IY	INFO	21

INFO3(IWAIT)=IZ	INFO	22
RETURN	INFO	23
20 CALL ERROR(3)	INFO	24
RETURN	INFO	25
END	INFO	26
* LIST8	SEEK	01
* LABEL	SEEK	02
SUBROUTINE SEEK(IX,IY,IZ)	SEEK	03
DIMENSION ICHRT(20000),ICOIN1(1000),ICOIN2(1000),IGEN(100,13)	SEEK	04
DIMENSION IPOWR(100),NLETR(100),ITBL(10,200),IWPG(10),ITALLY(100)	SEEK	05
DIMENSION INFO1(50),INFO2(50),INFO3(50)	SEEK	06
COMMON NWORD,NORD,NUM,NGENS,MAX,MAXROW,NDEF,IWORK,IWAIT,IR	SEEK	07
COMMON LPWMAX,LWPGMX,ICTMAX,IWMAX,ICMAX,IN,NOUT,ITR,NJOB,NJOBS	SEEK	08
COMMON ISCAN,ILOOK,ITALLY,INFO1,INFO2,INFO3,IWPG,ITBL	SEEK	09
COMMON NLETR,IPOWR,IGEN,ICOIN1,ICOIN2,ICHRT	SEEK	10
XLCTF(1,J)=2*NGENS*(I-1)+2*(XABSF(J)-1)+(3-XSIGNF(1,J))/2	SEEK	11
GO TO (3000,3000,3000,1000),ITR	SEEK	12
1000 WRITE OUTPUT TAPE NOUT,2000,IX,IY,IZ	SEEK	13
2000 FORMAT(6H SEEK(,I5,1H,,I5,1H,,I5,1H))	SEEK	14
3000 IF(IX) 10,20,20	SEEK	15
10 IZ=XLCTF(-IX,IY)	SEEK	16
RETURN	SEEK	17
20 IZ=XLCTF(IX,IY)	SEEK	18
IF(ICHRT(IZ)) 30,40,30	SEEK	19
30 IX=ICHRT(IZ)	SEEK	20
IZ=1	SEEK	21
RETURN	SEEK	22
40 IZ=2	SEEK	23
RETURN	SEEK	24
END	SEEK	25
	COINC	
* LIST8	COINC	01
* LABEL	COINC	02
SUBROUTINE COINC(INX,INY,INK)	COINC	03
DIMENSION ICHRT(20000),ICOIN1(1000),ICOIN2(1000),IGEN(100,13)	COINC	04
DIMENSION IPOWR(100),NLETR(100),ITBL(10,200),IWPG(10),ITALLY(100)	COINC	05
DIMENSION INFO1(50),INFO2(50),INFO3(50)	COINC	06
COMMON NWORD,NORD,NUM,NGENS,MAX,MAXROW,NDEF,IWORK,IWAIT,IR	COINC	07
COMMON LPWMAX,LWPGMX,ICTMAX,IWMAX,ICMAX,IN,NOUT,ITR,NJOB,NJOBS	COINC	08
COMMON ISCAN,ILOOK,ITALLY,INFO1,INFO2,INFO3,IWPG,ITBL	COINC	09
COMMON NLETR,IPOWR,IGEN,ICOIN1,ICOIN2,ICHRT	COINC	10
IF(INK) 25,25,10	COINC	11
10 GO TO (20,1000,1000,1000),ITR	COINC	12
1000 WRITE OUTPUT TAPE NOUT,2000,INX,INY,INK	COINC	13
2000 FORMAT(7H COINC(,I5,1H,,I5,1H,,I5,1H))	COINC	14
20 ITALLY(INK)=ITALLY(INK)+1	COINC	15
N=0	COINC	16
GO TO 30	COINC	17
25 N=-1	COINC	18
30 CALL ENTER(INX,INY,N)	COINC	19
40 IX=ICOIN1(1)	COINC	20
IY=ICOIN2(1)	COINC	21
IF(ILOOK-IY) 60,50,60	COINC	22

50	ILOOK=IX	COINC 23
60	IF(IWAIT) 120,120,70	COINC 24
70	DO 110 I=1,IWAIT	COINC 25
	IF(INFO1(I)-IY) 90,80,90	COINC 26
80	INFO1(I)=IX	COINC 27
90	IF(INFO3(I)-IY) 110,100,110	COINC 28
100	INFO3(I)=IX	COINC 29
110	CONTINUE	COINC 30
120	CALL SEEK(-IX,1,II)	COINC 31
	CALL SEEK(-IY,1,JJ)	COINC 32
	DO 210 K=1,NGENS	COINC 33
	DO 210 L=1,2	COINC 34
	M=2*K+L-3	COINC 35
	I=II+M	COINC 36
	J=JJ+M	COINC 37
	IF(ICHRT(J)) 130,210,130	COINC 38
130	IF(ICHRT(J)-IY) 150,140,150	COINC 39
140	ICHRT(J)=IX	COINC 40
	GO TO 155	COINC 41
150	CALL SEEK(-ICHRT(J),XSIGNF(K,2*L-3),M)	COINC 42
	ICHRT(M)=0	COINC 43
155	IF(ICHRT(I)) 160,190,160	COINC 44
160	IF(ICHRT(I)-IY) 180,170,180	COINC 45
170	ICHRT(I)=IX	COINC 46
180	CALL ENTER(ICHRT(I),ICHRT(J),N)	COINC 47
	GO TO 200	COINC 48
190	ICHRT(I)=ICHRT(J)	COINC 49
200	CALL SEEK(-ICHRT(I),XSIGNF(K,2*L-3),M)	COINC 50
	IF(ICHRT(M)) 206,203,206	COINC 51
203	ICHRT(M)=IX	COINC 52
206	ICHRT(J)=0	COINC 53
210	CONTINUE	COINC 54
	IF(INK) 250,250,220	COINC 55
220	IF(N-1) 250,250,230	COINC 56
230	N=N-1	COINC 57
	DO 240 I=1,N	COINC 58
	ICOIN1(I)=ICOIN1(I+1)	COINC 59
240	ICOIN2(I)=ICOIN2(I+1)	COINC 60
	ICOIN1(N+1)=0	COINC 61
	ICOIN2(N+1)=0	COINC 62
	GO TO 40	COINC 63
250	N=0	COINC 64
	RETURN	COINC 65
	END	COINC 66
*	LIST8	ENTER
*	LABEL	ENTER 01
	SUBROUTINE ENTER(INX,INY,N)	ENTER 02
	DIMENSION ICHRT(20000),ICOIN1(1000),ICOIN2(1000),IGEN(100,13)	ENTER 03
	DIMENSION IPOWR(100),NLETR(100),ITBL(10,200),IWPG(10),ITALLY(100)	ENTER 04
	DIMENSION INFO1(50),INFO2(50),INFO3(50)	ENTER 05
	COMMON NWORD,NORD,NUM,NGENS,MAX,MAXROW,NDEF,IWORK,IWAIT,IR	ENTER 06
	COMMON LPWMAX,LWPGMX,ICTMAX,IWMAX,ICMAX,IN,NOUT,ITR,NJOB,NJOBS	ENTER 07
	COMMON ISCAN,ILOOK,ITALLY,INFO1,INFO2,INFO3,IWPG,ITBL	ENTER 08
		ENTER 09

COMMON NLETR, IPOWR, IGEN, ICOIN1, ICOIN2, ICHRT	ENTER 10
IF(INX-INY) 10,170,20	ENTER 11
10 IX=INX	ENTER 12
IY=INY	ENTER 13
GO TO 30	ENTER 14
20 IX=INY	ENTER 15
IY=INX	ENTER 16
30 IF(N) 240,220,40	ENTER 17
40 IF(IY-ICOIN2(1)) 100,50,80	ENTER 18
50 IF(IX-ICOIN1(1)) 70,170,60	ENTER 19
60 IY=IX	ENTER 20
IX=ICOIN1(1)	ENTER 21
GO TO 100	ENTER 22
70 IY=ICOIN1(1)	ENTER 23
GO TO 100	ENTER 24
80 IF(IX-ICOIN2(1)) 100,90,100	ENTER 25
90 IX=ICOIN1(1)	ENTER 26
100 IF(N-1) 220,220,110	ENTER 27
110 DO 150 I=2,N	ENTER 28
IF(ICOIN2(I)-IY) 180,120,150	ENTER 29
120 IF(ICOIN1(I)-IX) 130,170,140	ENTER 30
130 IY=IX	ENTER 31
IX=ICOIN1(I)	ENTER 32
ICOIN1(I)=IY	ENTER 33
GO TO 110	ENTER 34
140 ICOIN1(I)=IX	ENTER 35
GO TO 110	ENTER 36
150 CONTINUE	ENTER 37
IF(N-ICMAX) 220,160,160	ENTER 38
160 CALL ERROR(4)	ENTER 39
170 RETURN	ENTER 40
180 IF(N-ICMAX) 190,160,160	ENTER 41
190 IF(N-IR) 196,196,193	ENTER 42
193 IR=N	ENTER 43
196 DO 200 J=1,N	ENTER 44
K=I+N-J	ENTER 45
ICOIN1(K+1)=ICOIN1(K)	ENTER 46
200 ICOIN2(K+1)=ICOIN2(K)	ENTER 47
N=N+1	ENTER 48
ICOIN1(1)=IX	ENTER 49
ICOIN2(1)=IY	ENTER 50
210 GO TO (170,170,1000,1000), ITR	ENTER 51
1000 WRITE OUTPUT TAPE NOUT,2000,IX,IY,N	ENTER 52
2000 FORMAT(7H ENTER(,I5,1H,,I5,1H,,I5,1H))	ENTER 53
GO TO 170	ENTER 54
220 IF(N-ICMAX) 223,160,160	ENTER 55
223 IF(N-IR) 230,230,226	ENTER 56
226 IR=N	ENTER 57
230 N=N+1	ENTER 58
ICOIN1(N)=IX	ENTER 59
ICOIN2(N)=IY	ENTER 60
GO TO 210	ENTER 61
240 ICOIN1(1)=IX	ENTER 62
ICOIN2(1)=IY	ENTER 63



N=1	ENTER 64
GO TO 170	ENTER 65
END	ENTER 66
	CLOSE
* LIST8	CLOSE 01
* LABEL	CLOSE 02
SUBROUTINE CLOSE	CLOSE 03
DIMENSION ICHRT(20000),ICOIN1(1000),ICOIN2(1000),IGEN(100,13)	CLOSE 04
DIMENSION IPOWR(100),NLETR(100),ITBL(10,200),IWPG(10),ITALLY(100)	CLOSE 05
DIMENSION INFO1(50),INFO2(50),INFO3(50)	CLOSE 06
COMMON NWORD,NORD,NUM,NGENS,MAX,MAXROW,NDEF,IWORK,IWAIT,IR	CLOSE 07
COMMON LPWMAX,LWPGMX,ICTMAX,IWMAX,ICMAX,IN,NOUT,ITR,NJOB,NJOBS	CLOSE 08
COMMON ISCAN,ILOOK,ITALLY,INFO1,INFO2,INFO3,IWPG,ITBL	CLOSE 09
COMMON NLETR,IPOWR,IGEN,ICOIN1,ICOIN2,ICHRT	CLOSE 10
L=2*NGENS	CLOSE 11
J=0	CLOSE 12
M=0	CLOSE 13
10 J=J+1	CLOSE 14
I2=L*J	CLOSE 15
I1=I2-L+1	CLOSE 16
DO 20 I=I1,I2	CLOSE 17
IF(ICHRT(I)) 30,20,30	CLOSE 18
20 CONTINUE	CLOSE 19
M=M+1	CLOSE 20
GO TO 50	CLOSE 21
30 IF(M) 40,50,40	CLOSE 22
40 CALL COINC(J,J-M,0)	CLOSE 23
50 IF(ISCAN-J) 70,60,70	CLOSE 24
60 ISCAN=J-M	CLOSE 25
70 IF(J-NDEF) 10,80,80	CLOSE 26
80 NDEF=J-M	CLOSE 27
RETURN	CLOSE 28
END	CLOSE 29
	FINISH
* LIST8	FINISH01
* LABEL	FINISH02
SUBROUTINE FINISH	FINISH03
DIMENSION ICHRT(20000),ICOIN1(1000),ICOIN2(1000),IGEN(100,13)	FINISH04
DIMENSION IPOWR(100),NLETR(100),ITBL(10,200),IWPG(10),ITALLY(100)	FINISH05
DIMENSION INFO1(50),INFO2(50),INFO3(50)	FINISH06
COMMON NWORD,NORD,NUM,NGENS,MAX,MAXROW,NDEF,IWORK,IWAIT,IR	FINISH07
COMMON LPWMAX,LWPGMX,ICTMAX,IWMAX,ICMAX,IN,NOUT,ITR,NJOB,NJOBS	FINISH08
COMMON ISCAN,ILOOK,ITALLY,INFO1,INFO2,INFO3,IWPG,ITBL	FINISH09
COMMON NLETR,IPOWR,IGEN,ICOIN1,ICOIN2,ICHRT	FINISH10
NORD=NDEF*NORD	FINISH11
WRITE OUTPUT TAPE NOUT,30,NORD,MAX,IR	FINISH12
WRITE OUTPUT TAPE NOUT,60	FINISH13
DO 9 I=1,NWORD	FINISH14
9 WRITE OUTPUT TAPE NOUT,70,I,ITALLY(I)	FINISH15
DO 10 I=1,NGENS	FINISH16
INFO1(I)=I	FINISH17
10 INFO3(I)=-I	FINISH18
WRITE OUTPUT TAPE NOUT,40,(INFO1(I),INFO3(I),I=1,NGENS)	FINISH19
DO 20 I=1,NDEF	FINISH20

J=2*(I-1)*NGENS+1	FINISH21
K=2*I*NGENS	FINISH22
20 WRITE OUTPUT TAPE NOUT,50,I,(ICHRT(L),L=J,K)	FINISH23
RETURN	FINISH24
30 FORMAT(26H)THE ORDER OF THE GROUP IS,18,5H MAX=,18,7H COINC=,15/)	FINISH25
40 FORMAT(6X,18I6/)	FINISH26
50 FORMAT(19I6)	FINISH27
60 FORMAT(10X,10HWORD TALLY/)	FINISH28
70 FORMAT(8X,2I5)	FINISH29
END	FINISH30
* LIST8	ERROR
* LABEL	ERROR 01
SUBROUTINE ERROR(I)	ERROR 02
DIMENSION ICHRT(20000),IC0IN1(1000),IC0IN2(1000),IGEN(100,13)	ERROR 03
DIMENSION IPOWR(100),NLETR(100),ITBL(10,200),IWPG(10),ITALLY(100)	ERROR 04
DIMENSION INFO1(50),INFO2(50),INFO3(50)	ERROR 05
COMMON NWORD,NORD,NUM,NGENS,MAX,MAXROW,NDEF,IWORK,IWAIT,IR	ERROR 06
COMMON LPWMAX,LWPGMX,ICTMAX,IWMAX,ICMAX,IN,NOUT,ITR,NJOB,NJOBS	ERROR 07
COMMON ISCAN,ILOOK,ITALLY,INFO1,INFO2,INFO3,IWPG,ITBL	ERROR 08
COMMON NLETR,IPOWR,IGEN,IC0IN1,IC0IN2,ICHRT	ERROR 09
GO TO (10,30,50,70),I	ERROR 10
10 WRITE OUTPUT TAPE NOUT,20	ERROR 11
GO TO 90	ERROR 12
20 FORMAT(34H)MAXIMUM WORDS/GENERATOR EXCEEDED.//)	ERROR 13
30 WRITE OUTPUT TAPE NOUT,40	ERROR 14
GO TO 90	ERROR 15
40 FORMAT(39H)MULTIPLICATION TABLE STORAGE EXCEEDED.//)	ERROR 16
50 WRITE OUTPUT TAPE NOUT,60	ERROR 17
GO TO 90	ERROR 18
60 FORMAT(44H)WAITING INFORMATION TABLE STORAGE EXCEEDED.//)	ERROR 19
70 WRITE OUTPUT TAPE NOUT,80	ERROR 20
GO TO 90	ERROR 21
80 FORMAT(37H)COINCIDENCE TABLE CAPACITY EXCEEDED.//)	ERROR 22
90 CALL CLOCK(6H ERROR,-1)	ERROR 23
CALL DUMP(ICHRT(20000),ICHRT(19800),2,ICHRT(200),NWORD,2)	ERROR 24
RETURN	ERROR 25
END	ERROR 26
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